

A history-oriented view of causality

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The causality relation yields a simple basis for computing event occurrences in dynamic systems using the rule:

occurs(F) if causes(E,F) & occurs(E)

Here *causes(E,F)* means that event *E* brings about or is directly responsible for event *F* and *occurs(E)* means that event *E* occurs. However, it can be quite difficult to compute the causality relation when statements about it refer to the past or the future of causing events. A natural example of such a statement, and its formalized version are:

An event of aircraft taking off from an airport towards a radar causes an event of detection by that radar in the future provided an event of its changing course does not occur in between, and provided in the past an event of its being instructed to fly low has not occurred.

*causes(flies_towards(aircraft(A),radar(R),T),in_range(aircraft(A),radar(R),T1)) if
detection_time(aircraft(A),radar(R),T,T1) &
not exists(X). X<T & occurs(told_to_fly_low(aircraft(A),X)) &
not exists(X).T<X<T1 & occurs(changes_course(aircraft(A),X)).*

Here *flies_towards(A,R,T)* denotes an event of object *A* taking off towards object *R* at time *T*. Similarly for *in_range(A,R,T)*, *told_to_fly_low(A,T)*, *changes_course(A,T)*. *detection_time(A,R,T,T1)* computes the time *T1* at which the flight path of *A*, starting at *T*, intersects coverage of *R*.

The difficulty arises from two sources. First, there is mutual recursion between this statement about *causes* and the above definition of *occurs*. This can be quite difficult to control. If occurring events are computed in increasing (decreasing) order of time how do we evaluate reference to the future (past)? Second, reference can also be to states. As time can be real-valued the number of distinct states, e.g. positions of a moving object, can be non-computably infinite. How do we keep a record of all these?

We present a formalism called DMOD and use it to propose solutions to these problems. More generally, it can be used to model and simulate dynamic systems. It is based upon the following fundamental assumption:

If we know the history of a system till time T, i.e. the sequence of events which occurred till T, we can compute the value of every state parameter at any point of time till T.

Let *F* be a condition about the system which takes *m+1* arguments, *m* ≥ 0, and the last argument ranges over real-valued time. Let *F* be defined for entities *a1,...,am,t*. Then the proposition *F(a1,...,am,t)* is said to

causal_connection(*E,HE,F,HEF*) if
E=flies_towards(*aircraft(X),radar(R),CT*) &
F=in_range(*aircraft(X),radar(R),FT*) &
causality_predicted(*E,HE,F*) &
prediction_unfalsified(*E,HE,F,HEF*).

causality_predicted(*E,HE,F,HEF*) if
E=flies_towards(*aircraft(X),radar(R),CT*) &
F=in_range(*aircraft(X),radar(R),FT*) &
detection_time(*aircraft(X),radar(R),[E/HE],FT*),
not member(*told_to_fly_low*(*aircraft(X),_,HE*)).

prediction_unfalsified(*E,HE,F,HEF*) if
E=flies_towards(*aircraft(X),radar(R),CT*) &
F=in_range(*aircraft(X),radar(R),FT*) &
not member(*changes_course*(*aircraft(X),_,HEF*)).

detection_time(*A,R,H,T*) is similar to that above although its third argument is the history upto the causing event. Relevant information about positions and velocities can be retrieved from it.

We now show how to use causality rules to compute which events occur in a system, i.e. its history. Let $S=E_0,E_1,E_2,\dots$ be a sequence of events sorted in increasing order of time-stamps. Then S is said to satisfy causal-soundness if for each $j, j \neq 0$, there exists $i, i < j$ such that *causal_connection*(E_i,HE_i,E_j,HE_iE_j) holds, where HE_i is the sequence E_0,\dots,E_{i-1} and HE_iE_j is the sequence E_{i+1},\dots,E_{j-1} . Note that the initial event E_0 is exempt from requiring a cause. Intuitively, a sequence is causally-sound if every event in it, except the first one, has a cause in it. Causal-soundness is similar to weak-causality property of de Kleer & Brown.

Let $S=E_0,E_1,E_2,\dots$ be a sequence of events sorted in increasing order of time-stamps. Then S is said to satisfy causal-completeness if for each i, j , if there is an event G such that *causal_connection*(E_i,HE_i,G,HE_iE_j) then G also appears in S , where HE_i is the sequence E_0,\dots,E_{i-1} and HE_iE_j is the sequence E_{i+1},\dots,E_{j-1} . Intuitively, a sequence is causally-complete if it contains all the events whose occurrence is required by the occurrence of the initial event and the causality rules.

Let E_0 be a special initial event for the system. Assume that E_0 has occurred. A history of the system is defined to be a sequence of events starting at E_0 which is both causally-sound and causally-complete. Intuitively, it contains all of the events whose occurrence is required by the occurrence of the initial event and the causality rules, and only these events. It can be shown that there is exactly one history in which each event possesses a distinct time-stamp. It can represent the history of a non-concurrent system.

To compute a history, let the initial event E_0 occur. Suppose the history $HF=E_0,E_1,\dots,E_m$ till a certain point of time has been computed. We need to compute the next event E_{m+1} . Let $S_m=\{F_1,F_2,\dots\}$ be the set of events where for each F_i , there exists an E_i such that *causal_connection*(E_i,HE_i,F_i,HE_iF_i) holds, where HE_i is the sequence E_0,E_1,\dots,E_{i-1} and HE_iF_i is the sequence E_{i+1},\dots,E_m . Then, as the history must be causally-complete, the next event, E_{m+1} must be the event in S_m with the least time-stamp. If E_{m+1} cannot be computed the algorithm halts. As there may be more than one event in S_m with least time-stamp the algorithm is non-deterministic. A different history would be computed for each choice of E_{m+1} signifying that the system is concurrent.